

# C-DAC Four Days Technology Workshop

*ON*

**Hybrid Computing – Coprocessors/Accelerators  
Power-Aware Computing – Performance of  
Applications Kernels**

**hyPACK-2013  
Mode 3 : Intel Xeon Phi Coprocessors**

**Lecture Topic :  
Intel Xeon-Phi Coprocessor : Cilk Plus**

*Venue : CMSD, UoHYD ; Date : October 15-18, 2013*

# An Overview of Xeon Phi – Prog. Shared Address Space Prog., - Cilk Plus

## Lecture Outline

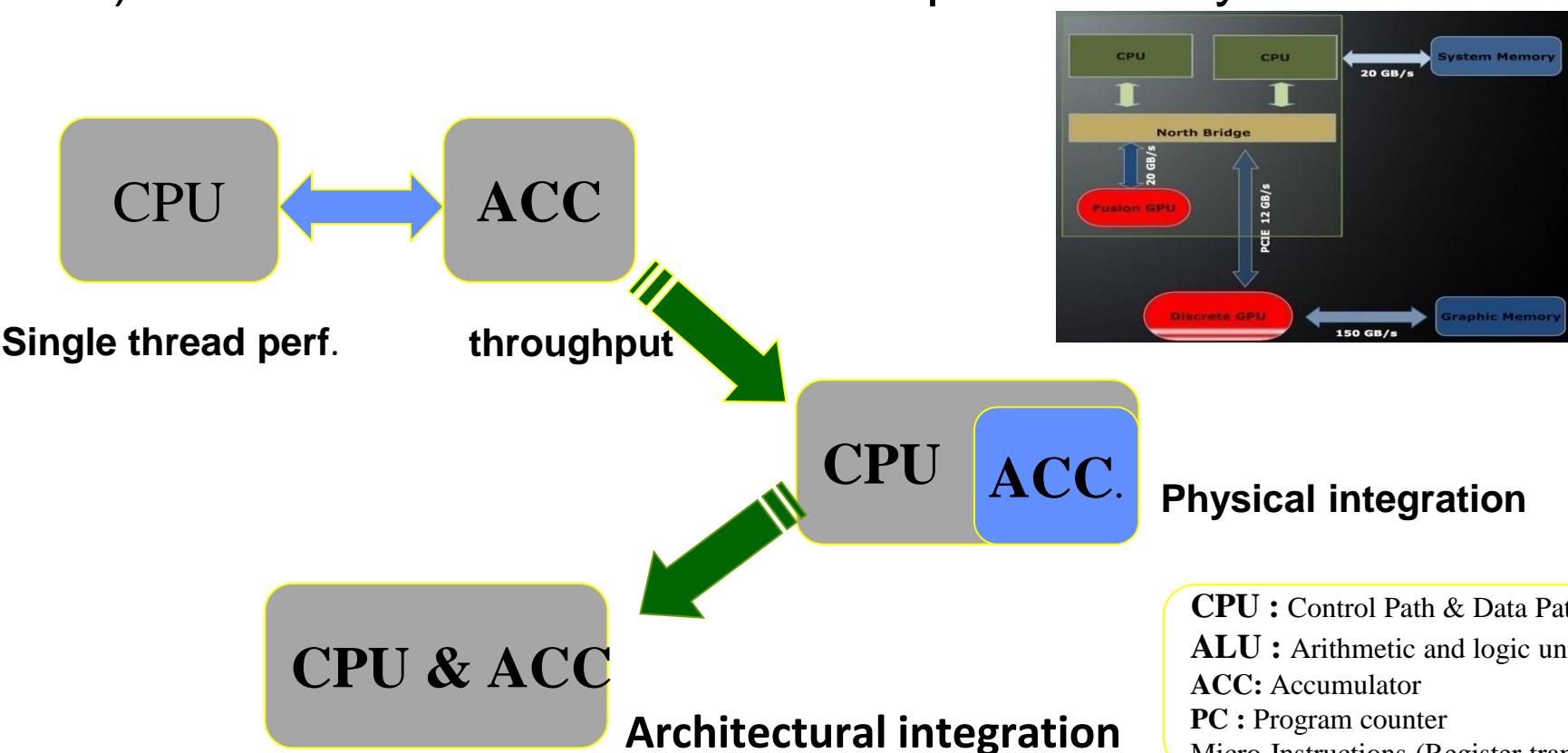
Following topics will be discussed

- ❖ Understanding of Xeon –Phi Architectures
- ❖ Programming on Xeon-Phi – Prog. -Cilk Plus
- ❖ Tuning & Performance – Software Threading

# Intel Xeon-Phi Shared Address Space Programming Cilk Plus

# Systems with Accelerators

A set (one or more) of very simple execution units that can perform few operations (with respect to standard CPU) with very high efficiency. When combined with full featured CPU (CISC or RISC) can accelerate the “nominal” speed of a system.

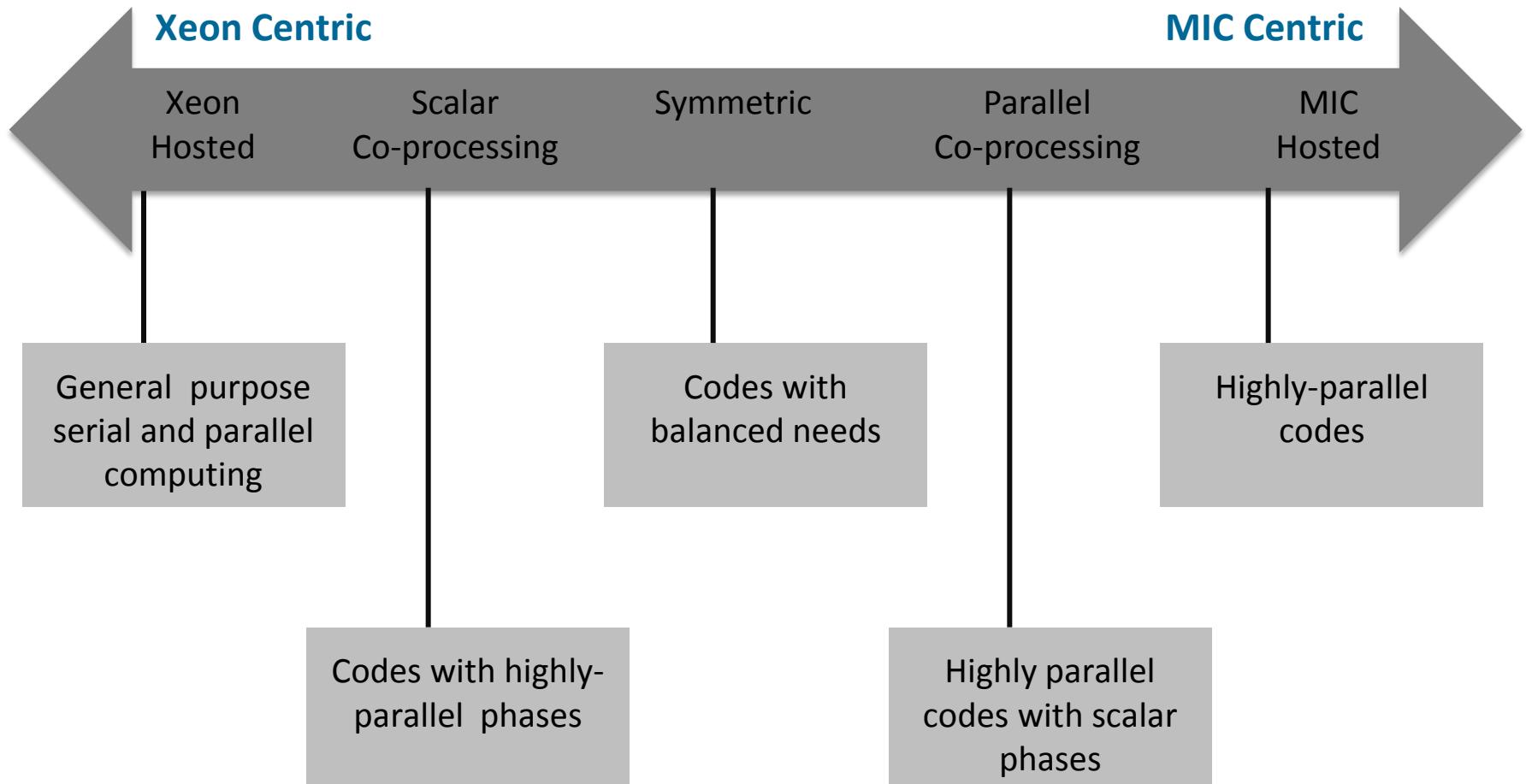


**CPU** : Control Path & Data Path  
**ALU** : Arithmetic and logic units  
**ACC** : Accumulator  
**PC** : Program counter  
Micro-Instructions (Register transfer)

**Source :** NVIDIA, AMD, SGI, Intel, IBM Alter, Xilinx References

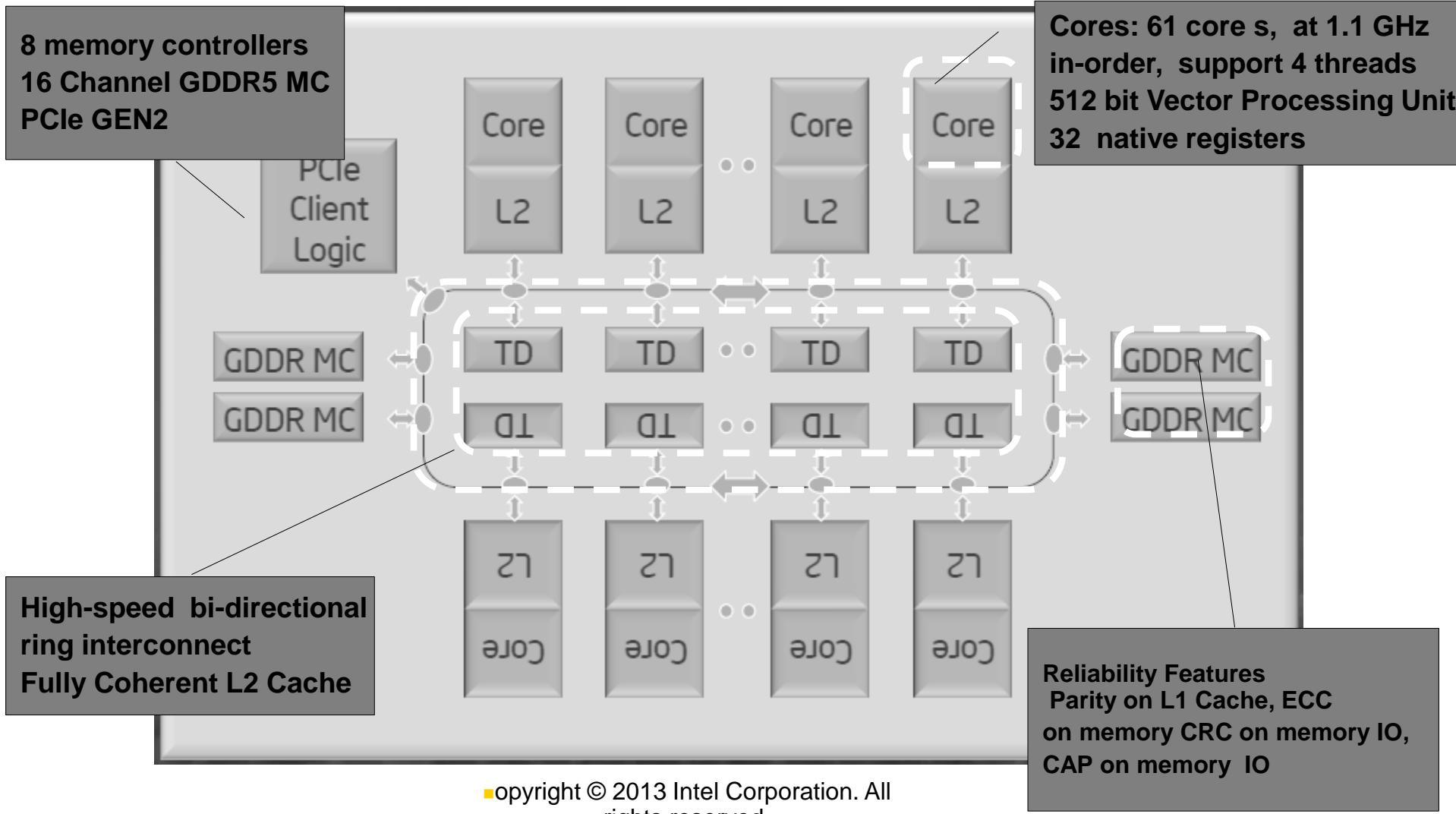
# MIC Architecture, System Overview

# Compute modes vision



Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel® Xeon Phi™ Architecture Overview



Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi Shared Address Space Programming Cilk Plus

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

- ❖ MIT Cilk – The original research project from MIT , culminating in Cilk-5.4.6. MIT Cilk was implemented as a source-to-source translator that converts Cilk code to C and then compiled the resulting C source.
  - Only supported C code with Cilk keywords
  - All parallel functions had to be marked with a cilk keyword
  - Cilk functions had to be spawned, not called

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

- ❖ **Cilk++** - Compilers developed by **Cilk Arts, Inc.** **Cilk** Arts licensed the Cilk technology from MIT.
  - Only supported C++ code
  - Used a non-standard calling convention, meaning you had to use a **cilk::context** to
  - call Cilk functions from C or C++
  - Cilk files used the **.cilk** extension
  - Released by Intel as unsupported software through the **WhatIf** site

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

- ❖ **Cilk++** - Intel Cilk Plus – Fully integrated into the Intel C/C++ compiler
  - Supports both C and C++
  - Uses standard calling conventions
  - Includes both task and data parallelism

# **Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus**

- ❖ **Why to use it ?**
  - Intel® Cilk™ Plus is the easiest, quickest way to harness the power of both multicore and vector processing.
- ❖ **What is it ?**
  - Intel Cilk Plus is an extension to the C and C++ languages to support data
- ❖ **Primary Features :**
  - HPC**
    - An efficient work-stealing scheduler provides nearly optimal scheduling of parallel tasks
    - Vector support unlocks the performance that's been hiding in your processors
    - Powerful hyperobjects allow for lock-free programming

# **Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus**

## **Primary Features :**

- ❖ **Easy to Learn**
  - Only 3 new keywords to implement task parallelism
  - Serial semantics make understanding and debugging the parallel program easier
  - Array Notations provide a natural way to express data parallelism
  
- ❖ **Easy to Use**
  - Automatic load balancing provides good behaviour in multi-programmed environments
  - Existing algorithms easily adapted for parallelism with minimal modification
  - Supports both C and C++ programmers

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Primary Features :

**Keywords :** Simple, powerful expression of task parallelism:

- **cilk\_for** - Parallelize for loops
- **cilk\_spawn** - Specifies that a function can execute inparallel with the remainder of the calling function
- **cilk\_sync** - Specifies that all spawned calls in a function must complete before execution continues

# Other Options for Parallelism: Intel® Cilk™ Plus

C/C++ extension for fine-grained task parallelism. 3 keywords:

## \_Cilk\_spawn

- ❖ Function call *may* be run in parallel with caller – up to the runtime

## \_Cilk\_sync

- ❖ Caller waits for all children to complete

## \_Cilk\_for

- ❖ Iterations are structured into a work queue
- ❖ Busy cores do not execute the loop
- ❖ Idle cores steal work items from the queue
- ❖ Countable loop Granularity is N/2, N/4, N/8, for trip count of N
- ❖ Intended use:
  - when iterations are not balanced, or
  - When overall load is not known at design time

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Primary Features :

**Keywords :** Simple, powerful expression of task parallelism:

- **Reducers:** Eliminate contention for shared variables among tasks by automatically creating views of them as needed and "reducing" them in a lock free manner.
- **Array Notation :** Data parallelism for arrays or sections of arrays.
- **Elemental Functions :** Define functions that can be vectorized when called from within an array notation expression or a **#pragma simd** loop
- **#pragma simd:** Specifies that a loop is to be vectorized

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi : Prog. Env. Cilk Plus

## Cilk Plus Keywords

cilk\_spawn and cilk\_sync:

Example of Fibonacci number calculator program

### ❖ Sequential

```
int fib(int n)
{
    if (n < 2)
        return n;
    int x = fib(n-1);
    int y = fib(n-2);
    return x + y;
}
```

### ❖ (With Cilk Plus Key Words)

```
int fib(int n)
{
    if (n < 2)
        return n;
    int x = cilk_spawn fib(n-1);
    int y = fib(n-2);
cilk_sync;
    return x + y;
}
```

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi : Programming Env.

## Offload Code Examples

### ❖ C/C+ Offload Pragma

```
#pragma offload target (mic)
#pragma omp parallel for reduction(+:pi)
for (i = 0; i<count; i++) {
    float t = (float) (i+0.5/count);
    pi += 4.0/(1.0t*t);
}
pi/ = count;
```

### ❖ C/C++ Offload Pragma

```
#pragma offload target(mic)
    in(transa, transb, N, alpha, beta) \
    in(A:length(matrix_elements)) \
    in(B:length(matrix_elements)) \
    inout(C:length(matrix_elements))
        sgemm(&transa, &transb, &N, &N, &N,
& alpha, A, &N, B, & N, &beta, C &N);
```

### ❖ Fortran Offload Directives

```
!dir$ omp offload target(mic)
 !$omp parallel do
    do i = 1, 10
        A(i) = B(i) * C(i)
    enddo
```

### ❖ C/C++ Language Extension

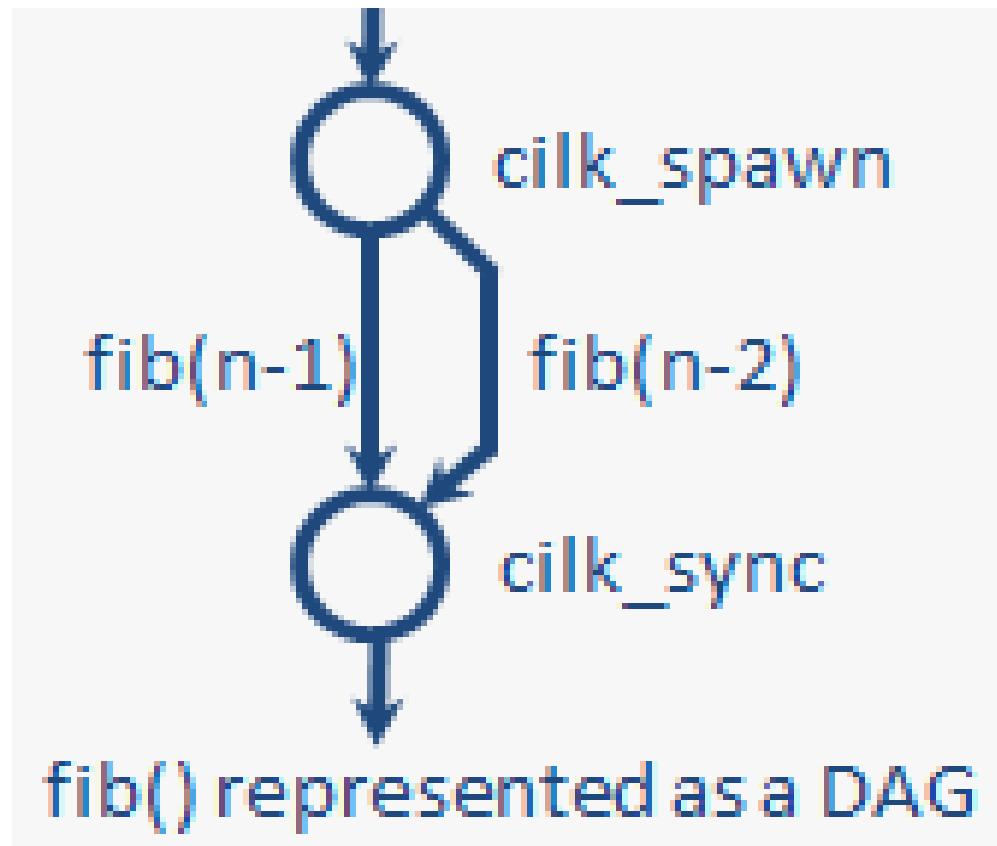
```
class_Cilk_Shared common {
    int data1;
    int *data2;
    class common *next;
    void process();
}
_Cilk_Shared class common obj1, obj2;
_Cilk_spawn _offload obj1.process();
_Cilk_spawn _offload obj2.process();
```

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi : Prog. Env. Cilk Plus

## Cilk Plus Keywords

cilk\_spawn and cilk\_sync:



<http://www.intel.com/>

# Intel Xeon-Phi : Prog. Env. Cilk Plus

## Cilk Plus Keywords

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Example of Fibonacci number calculator program

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```
int fib(int n)
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```

### ❖ (With Cilk Plus Key Words)

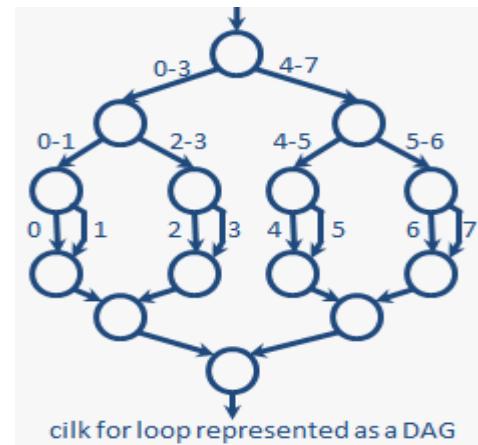
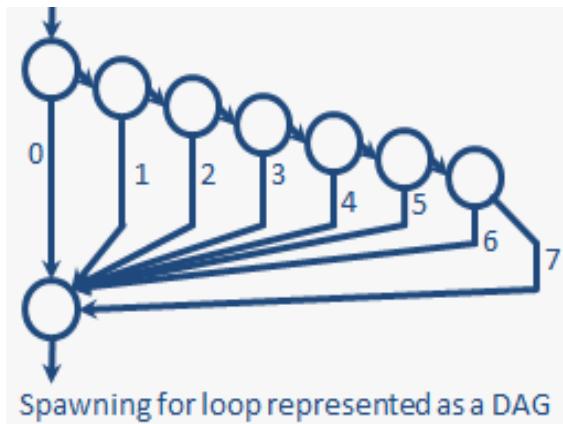
```
int fib(int n)
{
    if (n < 2)
        return n;
    int x = cilk_spawn fib(n-1);
    int y = fib(n-2);
cilk_sync;
    return x + y;
}
```

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi : Prog. Env. Cilk Plus

## Cilk Plus Keywords

Advantage of `cilk_for` over `cilk_spawn`



### **cilk\_spawn code**

```
for (int i = 0; i < 8; ++i)
{
    cilk_spawn do_work(i);
}
cilk_sync;
```

### **cilk\_for code**

```
cilk_for (int i = 0; i < 8; ++i)
{
    do_work(i);
}
```

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Cilk Plus Keywords

### ❖ Features of **cilk\_spawn**:

- **cilk\_spawn** permits parallelism. It does not command it. **cilk\_spawn** does not create a thread. It allows the runtime to steal the **continuation** to execute in another worker thread.
- A **strand** is a sequence of instructions that starts or ends on a statement which will change the parallelism.
- Permitting parallelism instead of commanding it is an aspect of the **serial semantics** of a deterministic **Intel Cilk Plus** application.
- It is always possible to run an **Intel Cilk Plus** application with a single worker, and it should give identical results to the **serialization** of that program

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Cilk Plus Reducers : The Cilk Plus Reducer Library :

### Lists

`reducer_list_append`: Creates a list by adding elements to the back.  
`reducer_listprepend`: Creates a list by adding elements to the front.

### Min/Max

`reducer_max`: Calculates the maximum value of a set of values.  
`reducer_max_index`: Calculates the maximum value and index of that value of a set of values.

`reducer_min`: Calculates the minimum value of a set of values.  
`reducer_min_index`: Calculates the minimum value and index of that value of a set of values.

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Cilk Plus Reducers : The Cilk Plus Reducer Library :

### Math Operators

`reducer_opadd`: Calculates the sum of a set of values.

### Bitwise Operators

- ❖ `reducer_opand`: Calculates the binary AND of a set of values.
- ❖ `reducer_opor`: Calculate the binary OR of a set of values.
- ❖ `reducer_opxor`: Calculate the binary XOR of a set of values.

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Cilk Plus Reducers : The Cilk Plus Reducer Library :

```
void locked_list_test()
{ mutex m;
  std::list<char>letters;
  // Build the list in parallel
  cilk_for(char ch = 'a'; ch <= 'z'; ch++)
  { simulated_work();
    m.lock();
    letters.push_back(ch);
    m.unlock(); }
  // Show the resulting list
  std::cout << "Letters from locked list: ";
  for(std::list<char>::iterator i = letters.begin(); i != letters.end(); i++)
  { std::cout << " " << *i;
  }std::cout << std::endl;}
```

Letters from locked list: y g n d t a w x e z q j o h b u f v c k i r p l m s

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Cilk Plus Reducers : The Cilk Plus Reducer Library :

```
void reducer_list_test()
{ cilk::reducer_list_append<char> letters_reducer;
  // Build the list in parallel
  cilk_for(char ch = 'a'; ch <= 'z'; ch++)
  { simulated_work();
    letters_reducer.push_back(ch);
  }
  // Fetch the result of the reducer as a standard STL list
  const std::list<char> &letters = letters_reducer.get_value();
  // Show the resulting list
  std::cout << "Letters from reducer_list:";
  for(std::list<char>::const_iterator i = letters.begin(); i != letters.end(); i++)
  { std::cout << " " << *i;
  }std::cout << std::endl;}
```

Letters from reducer\_list: a b c d e f g h i j k l m n o p q r s t u v w x y z

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel® Cilk™ Plus Array Notation

- ❖ C/C++ Language extension supported by the Intel® Compiler

- ❖ Based on the concept of array-section notation:

```
<array>[<low_bound> : <len> : <stride>] [<low_bound> : <len> :  
<stride>]...
```

- ❖ C/C++ Operators / Function Calls

- `d[:] = a[:] + (b[:] * c[:])`
- `b[:] = exp(a[:]);` // Call `exp()` on each element of `a[]`

- ❖ Reductions combine array section elements to generate a scalar result

- Nine built-in reduction functions supporting basic C data-types:
  - add, mul, max, max\_ind, min, min\_ind, all\_zero, all\_non\_zero, any\_nonzero
- Supports user-defined reduction function
  - Built-in reductions provide best performance

```
float a[10];  
..  
= c[][][5];
```

```
float a[10];
```

```
..
```

```
= a[:];
```

```
float a[10];
```

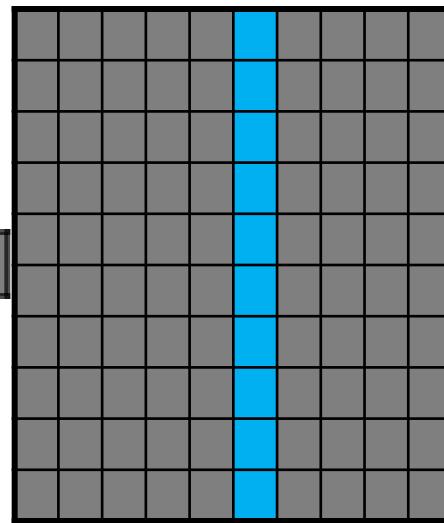
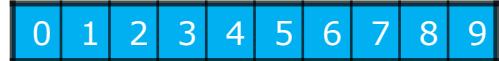
```
..
```

```
= a[2:6];
```

```
float a[10];
```

```
..
```

```
= d[0:3:2];
```



## Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

### Cilk Plus Array Notation

- ❖ Intel Cilk Plus includes extensions to C and C++ that allows for parallel operations on arrays.
- ❖ The intent is to allow users to express high-level vector parallel array operations. –
- ❖ This helps the compiler to effectively vectorize the code.
- ❖ Array notation can be used for both static and dynamic arrays.
- ❖ The extension has parallel semantics that are well suited for per-element operations that have no implied ordering and are intended to execute in data-parallel instructions.

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Cilk Plus Array Notation

A new operator [:] delineates an array section:

**array-expression [lower-bound : length : stride]**

- ❖ Length is used instead of upper bound as C and C++ arrays are declared with a given length.
- ❖ The three elements of the array notation can be any integer expression. The user is responsible for keeping the section within the bounds of the array.
- ❖ Each **argument to [:]** may be omitted if the default value is sufficient.
  - The default lower-bound is 0 .
  - The default length is the length of the array. If the length of the array is not known, length must be specified.
  - The default stride is 1. If the stride is defaulted, the second ":" may be omitted.

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## User-mandated Vectorization(pragma simd)

- ❖ SIMD (Single Instruction, Multiple Data) vectorization uses the **#pragma simd** pragma to enforce loop vectorization.
- ❖ Consider an example in C++ where the function **add\_floats()** uses too many unknown pointers, preventing automatic vectorization. You can give a data-dependence assertion using the **auto-vectorization** hint via **#pragma ivdep** and let the compiler decide whether the **auto-vectorization** optimization should be applied to the loop. Or you can now enforce vectorization of this loop by using **#pragma simd**.

```
void add_floats(float *a, float *b, float *c,
                float *d, float *e, int n)
{
    int i;
    #pragma simd
    for (i=0; i<n; i++) {
        a[i] = a[i]+ b[i]+c[i]+d[i] + e[i];
    }
}
```

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Cilk Plus : User-mandated Vectorization(pragma SIMD)

- ❖ The one big difference between using the SIMD pragma and **auto-vectorization** hints is that with the SIMD pragma, the compiler generates a warning when it is unable to vectorize the loop. With **auto-vectorization** hints, actual vectorization is still under the discretion of the compiler, even when you use the **#pragma vector** always hint.
- ❖ If a **#pragma SIMD** annotated loop is not vectorized by the compiler, the loop holds its serial semantics.
- ❖ By default “**#pragma SIMD**” is set to “**noassert**”.and compiler will issue a warning if the loop fails to **vectorize**.
- ❖ To direct the compiler to assert an error when the **#pragma SIMD** annotated loop fails to vectorize , add the “**assert**” clause to the **#pragma SIMD**

# Intel® Cilk™ Plus Technology: Elemental Function

- ❖ Allow you to define data operations using scalar syntax
- ❖ Compiler apply the operation to data arrays in parallel, utilizing both SIMD parallelism and core parallelism

## Programmer

1. Writes a standard C/C++ scalar syntax
2. Annotate it with `__declspec(vector)`
3. Use one of the parallel syntax choices to invoke the function

## Build with Intel Cilk Plus Compiler

1. Generates vector code with SIMD Instr.
2. Invokes the function iteratively, until all elements are processed
3. Execute on a single core, or use the task scheduler, execute on multicores

```
    __declspec (vector)
double BlackScholesCall(double S,
                        double K,
                        double T)
{
    double d1, d2, sqrtT = sqrt(T);
d1 = (log(S/K)+R*T) / (V*sqrtT)+0.5*V*sqrtT;
    d2 = d1-(V*sqrtT);
    return S*CND(d1) - K*exp(-R*T)*CND(d2);
}
```

```
Cilk_for (int i=0; i < NUM_OPTIONS; i++)
    call[i] = BlackScholesCall(SList[i],
                               KList[i],
                               TList[i]);
```

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Elemental Functions

- ❖ An elemental function is a regular function, which can be invoked either on scalar arguments or on array elements in parallel. You define an elemental function by adding

“`__declspec(vector)`” (on Windows\*) and  
“`__attribute__((vector))`” (on Linux\*) before

- ❖ the function signature:

```
__declspec (vector)
double ef_add (double x, double y) {return x + y;}
```

When you declare a function as elemental the compiler generates a short vector form of the function, which can perform your function's operation on multiple arguments in a single invocation.

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

# Intel Xeon-Phi - for Parallelism: Intel® Cilk™ Plus

## Elemental Functions

The vector form of the function can be invoked in parallel contexts in the following ways:

1. From a for-loop. It gets auto-vectorized; a loop that only has a call to an elemental function is always vectorizable, but the **auto-vectorizer** is allowed to apply performance heuristics and decide not to vectorize the function.
2. From a for-loop with **pragma simd**. If the elemental function is called from a loop with "**pragma simd**", the compiler no longer does any performance heuristics, and is guaranteed to call the vector version of the function.
3. From a **cilk\_for**
4. From an array notation syntax. .

## Summary: Tricks for Performance

- ❖ Use asynchronous data transfer and double buffering offloads to overlap the communication with the computation
- ❖ Optimizing memory use on Intel MIC architecture target relies on understanding access patterns
- ❖ Many old tricks still apply: peeling, collapsing, unrolling, vectorization can all benefit performance

Source : References & Intel Xeon-Phi; <http://www.intel.com/>

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# An Overview of Multi-Core Processors

## Conclusions

- ❖ An Overview of Xeon-Phi Architectures, Programming on based on Shared Address Space Platforms – Cilk Plus, Performance of Software threading are discussed.

Thank You  
*Any questions ?*